Rapid VM Synchronization with I/O Emulation Logging-Replay

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Motivation

- VM replication including backend storage is important
 - enables use of shared-nothing architecture
 - enhances VM mobility
- However, it heavily degrades VM performance



Our goal and approach

- Goal
 - Developing a rapid VM (including backend storage) synchronization mechanism
- Approach
 - logging-replay: The secondary VM replays the disk I/O events produced by the primary VM.



How to replicate VM and disk data?



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How to replicate VM and disk data?



- Transfer I/O logs instead of dirty blocks to the secondary VM – Transfer VM info(dirty pages and device info) to the secondary VM as usual
- The secondary VM only replays disk write events using I/O logs

- Needs to be the same VM memory as the primary VM's when secondary VM replays disk write events
 - Secondary VM can't make the same VM memory as the primary VM's
- Should we sync VMs when primary VM runs disk write events?

Specific event-driven VM synchronization

Disk write using two kinds of data

 dirty pages and non-dirty pages



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Implementation

• Kemari: VM synchronization mechanism for achieving fault tolerance

Needs shared disks

 Kemari with logging-replay can become a shared-nothing architecture



Architecture based on QEMU/KVM



- I/O-logging: saves the disk write events' logs
- I/O-replay: replays the disk write events using I/O logs

I/O-logging: saves the disk write events' logs



- Saves I/O logs at block layer in QEMU
 - applicable to many device models
- Need to save the following information to replay disk I/O events:
 - device name
 - \Rightarrow to get BlockDriverState structure
 - memory page address
 - \Rightarrow to get disk write data
 - (memory data transferred already)
 - num and location of disk sectors
 - \Rightarrow to write data

I/O-Replay: replays the disk write events using I/O logs



- Secondary VM requests I/O events using I/O logs from blocklayer
- I/O-Replay's callback function is invoked when disk write is finished
- I/O-replay runs at the block-layer
 - device info is not modified

Experimentation

- Experiment items
 - Traffic needed to synchronize VMs and disks
 - Performance of the primary VM (File I/O) using IOzone
- Techniques implemented in Kemari:
 - logging-replay
 - dirty-block-copy
 - Transferring dirty blocks to secondary VM
- Experimental environment
 - Hardware specs:
 - CPU: Quad-core Intel Xeon 2.6GHz X 2
 - Network: 1 GbE
 - VM specs:
 - KVM: Linux 2.6.33
 - QEMU: qemu-0.14.0
 - Guest OS: Debian lenny w/ virtio-blk
 - Memory: 1 GB

Traffic



 Logging-replay reduces the traffic neede synchronize only disks by 99%

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Synchronized NW traffic (1 Gbps)



Feature of event-driven synchronization

File I/O for both bandwidths



 Good performance obtained at both bandwidths, especially at low bandwidth

Summary

- A rapid VM synchronization mechanism is being developed.
- Logging-replay was implemented in Kemari.
 - Traffic/sync reduced by 50%
 - Throughput increased by 160%

Future work

- Addressing the problem of disk failure propagation
 - Disk read errors propagate to secondary VM when primary VM's disk crashes.
 - In other failure cases, we can verify disk consistency.
- Implementing memory replication with logging-replay
 - Enabling secondary VM to replay disk read events produced by the primary VM
- Implementing live migration with loggingreplay to improve performance