

Linux Storage Stack for the Cloud

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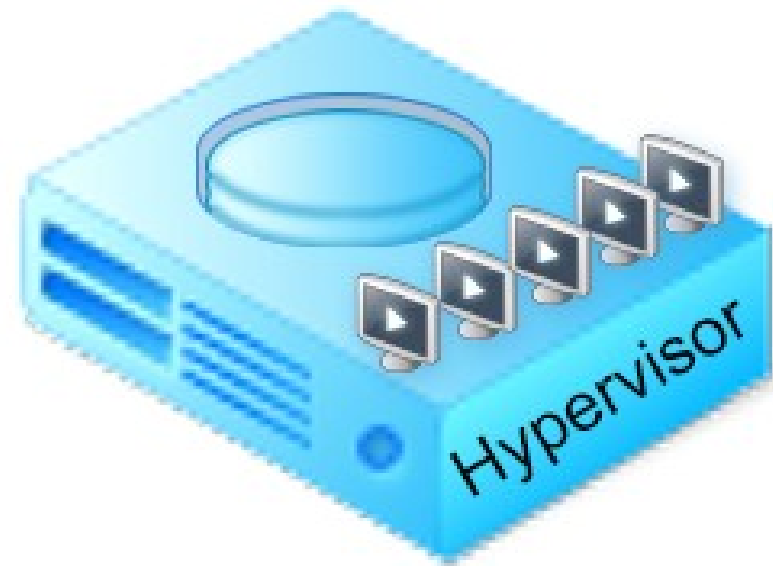
Who am I



- Storage virtualization – Why, What and How?
- Challenges & Solutions in the enterprise
- oVirt Design and Implementation
- Q&A

Why storage virtualization?

- Limited physical disk interfaces
- Fixed size
- Can't join disks
 - Performance
 - Storage array limitations
 - Multiple arrays



Create virtual devices with disk behavior

- Partition table
- Storage arrays
- LVM

- **Space flexibility**
- **Create devices ‘on the fly’**
- **Snapshots**

A virtual disk for a vm

- One image is worth many volumes
- Volume:
 - YABS – Yet Another Block Sequence
- Volume types:
 - File
 - Block

What is
the problem?

Multiple **data centers**

x

hundreds of **hosts**

x

hundreds of **VMs**

x

multiple **disks**

x

potentially dozens of **snapshots**

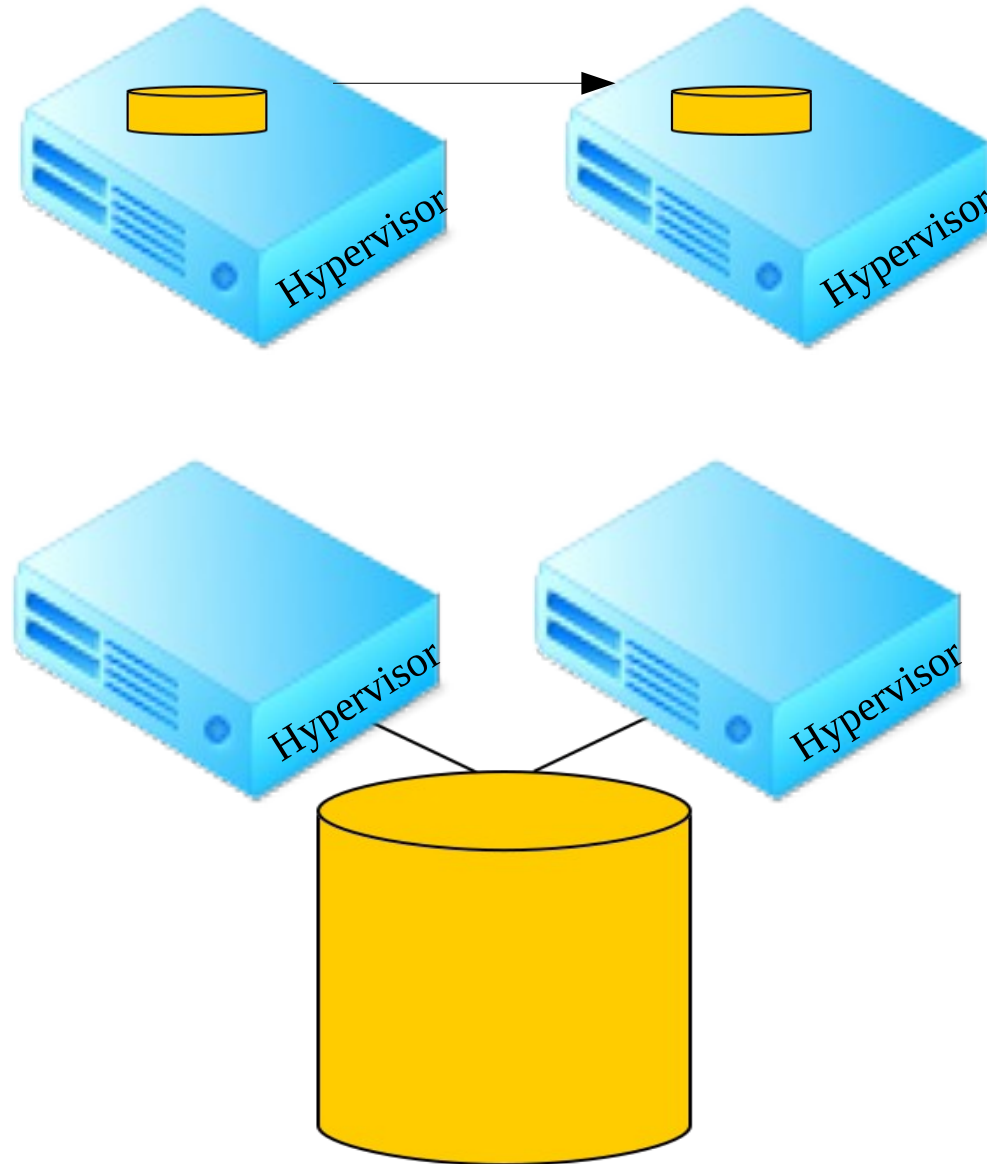
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VERY BIG HEADACHE



- **Host independent VMs**
- **Quantity of volumes**
- **Size of storage**

Host independent VMs



- **Host independent VMs**
 - Shared storage
- **Quantity**
 - Creation on the fly
 - Templates
 - Centralized DB
- **Size**
 - Over-commitment
 - Thin provisioning
 - Templates (Shared data, same OS)

oVirt Implementation

- Use **qcow2**
- file and block volumes
- provides COW volumes
- Thin volumes

File

- **Quantity**
 - create and manage files using the file system
 - “Unlimited”
- **Size**
 - Dynamic sizing
 - Sparse files
- **Shared storage**
 - NAS
 - Synchronizing access

Block

- **Quantity**
 - How do we create a block device?
 - How many block devices are supported?
- **Size**
 - How can we resize a block volume?
 - Is thin provisioning possible?
- **Shared storage**

Using remote storage. But...



- Different storage vendors, models
- No standard interface

Why Block?

- File system performance overhead
- Customer requirements

- Initiator, Target, LUN = **GUID**
- **Transport** for the SCSI commands
 - FC
 - iSCSI
- **Redundancy**
 - Multiple targets for the same LUN
 - How can we tell if it's the same LUN?

- **Using Multipath**
 - Query the storage to obtain the GUID
 - A new GUID is mapped through **device-mapper**
 - Use rules to choose the preferred path for the device
- Fail fast
- Pause VM
- I/O failure never reaches guest OS
- Auto resume

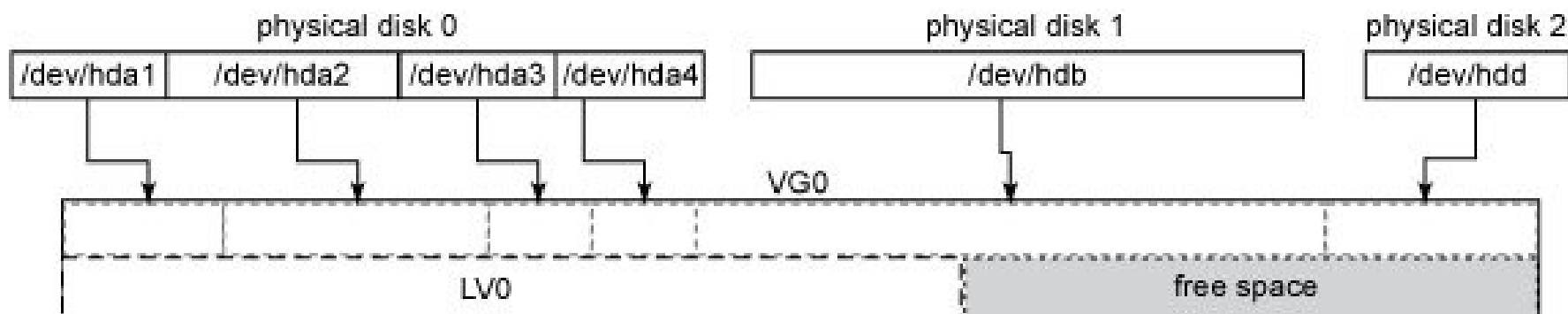
Why device-mapper?



- mapping block devices onto virtual block devices
- Used by multiple Linux storage stack components
- Multipath, RAID, LVM, crypt, etc...

Creating and managing block images

- LVM provides a unified interface
- Volume is implemented as an LV
- Easy provisioning: `lvcreate`, `lvremove`
- Thin provisioning: `lvextend`



* <http://www.markus-gattol.name/ws/lvm.html>

**Very
specialized
use of
LVM**

- No use of LVM native thin provisioning
- LV initial size – 1GB
- Extend LV when:
 - VM paused due to ENOSPC
 - High watermark (monitoring qemu) identified

Need a clustered solution



- create, remove, extend are VG MD writes
- Simultaneous writes will cause MD corruption
- cLVM did not scale
- No synchronization mechanisms

- **Hybrid mode and compartmentalization**
 - Runtime config, separate for vdsmd
 - to avoid affecting anything else on the host
 - Allow admin to make changes outside of vdsmd
 - LVM short filters
 - Speed up operations (by default LVM scans all devices)
 - Compartmentalize problems
 - Avoid accessing host 'owned' devices
- **Activate / deactivate**
 - Keep number of devices lower
 - Avoid refresh

- LVM MDA per PV by default
 - Problems
 - In clustered environment with more than 1 PV **will** cause corruption
 - Requires update of multiple areas to commit transaction
 - Solution
 - only 1 active MDA
- oVirt MD as LV and VG tags
- Lock type 4 (patches upstream)

- **Storage Pool Manager**
- **A role** assigned to one host
- Can be migrated to any host in a data center
- Creation, deletion and manipulation of volumes
- **Single meta data writer**

- Cluster membership based on
 - Light-weight leases for storage-centric coordination (Chockler and Malkhi 2004)
- Single recoverable leader
- Primitives: lease and renew
- Uniform
- Simple and efficient

- Cluster membership, like SPM, based on
 - Light-weight leases for storage-centric coordination (Chockler and Malkhi 2004)
- Leases based on
 - Disk Paxos (modified for leases)

- Storage virtualization
- oVirt implementation
- oVirt snapshot
- File implementation
- Block implementation
- Multipath
- Device-mapper
- LVM
- SPM



THANK YOU !

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