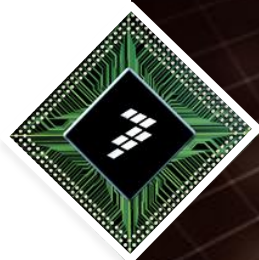




Efficient Guest Agnostic Virtualization With Embedded Power Architecture[®] KVM

Varun Sethi
Freescale Semiconductor



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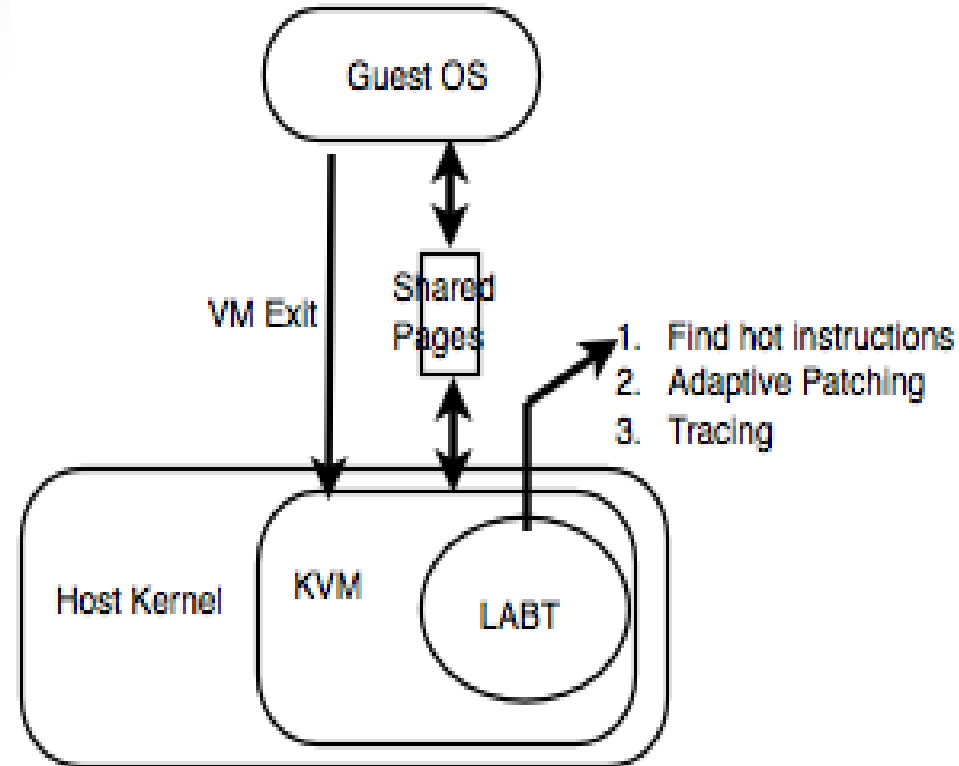


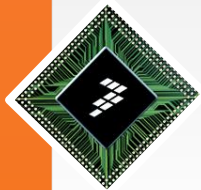
Motivation

- Growing interest in virtualization using KVM on embedded Power Architecture platforms
- Requirement to run various customer specific operating systems with embedded Power Architecture KVM
- Possible to run unmodified guest on embedded Power Architecture KVM – But this comes at a significant cost associated with VM exits
 - *Problem severe for cores without virtualization assists*
- Possible to “Binary translate” guest privileged instructions from the host side.
 - Continue to run an unmodified guest

Adaptive Binary Translation

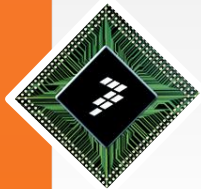
- Dynamically infer the instructions that cause a large number of VM exits
- Binary Translate these instructions to a faster emulation code .
- Instruction translation maintained on a shared memory region
- “Memory tracing” (remove R/W permissions to shared page) implemented to control guest access to modified pages
 - Guest access to page generates a DSI exception



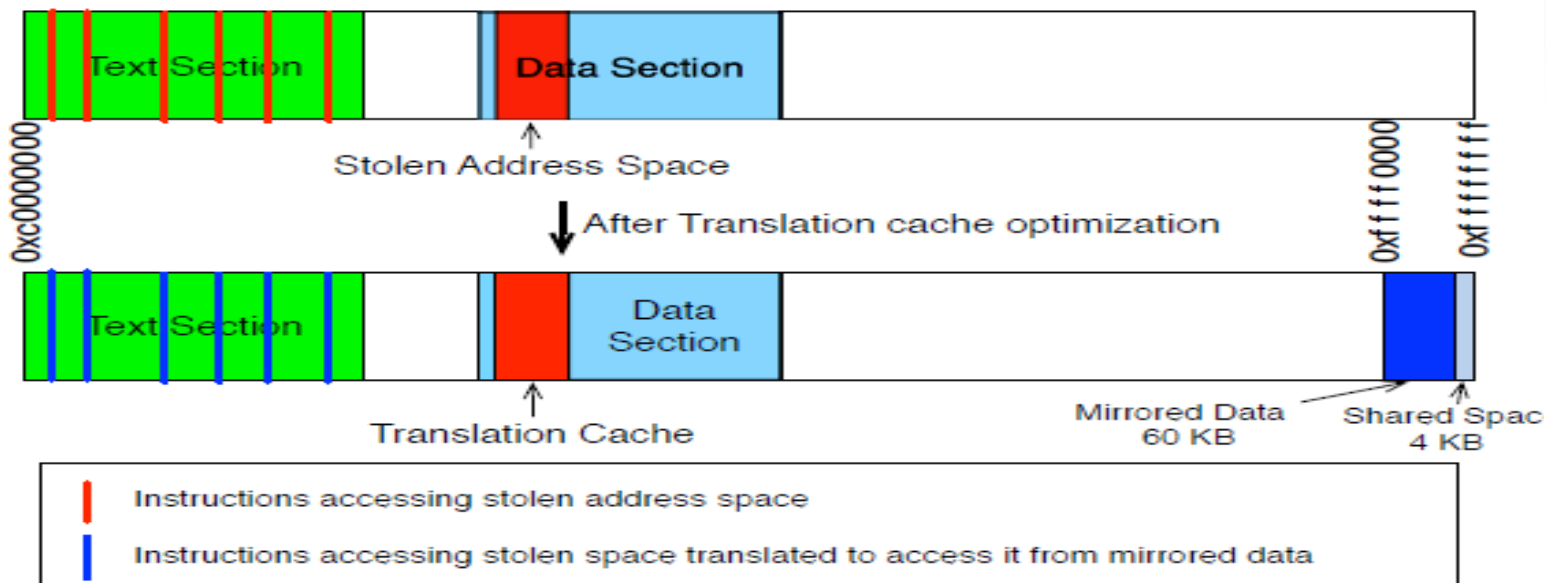


Design Challenges

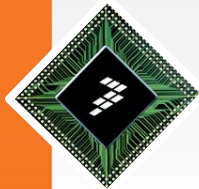
- Handling complex instruction translations
 - Multi line patching could be complex
 - Translation cache placement issue
- Minimizing performance overhead (excessive DSIs) due to Memory tracing
 - Use of huge pages (TLB1)
 - Self referential code
 - Access to sys call table and exception prolog in case of Linux
 - Self modifying code
 - Code from the modified page trying to modify (Write) code on the same page



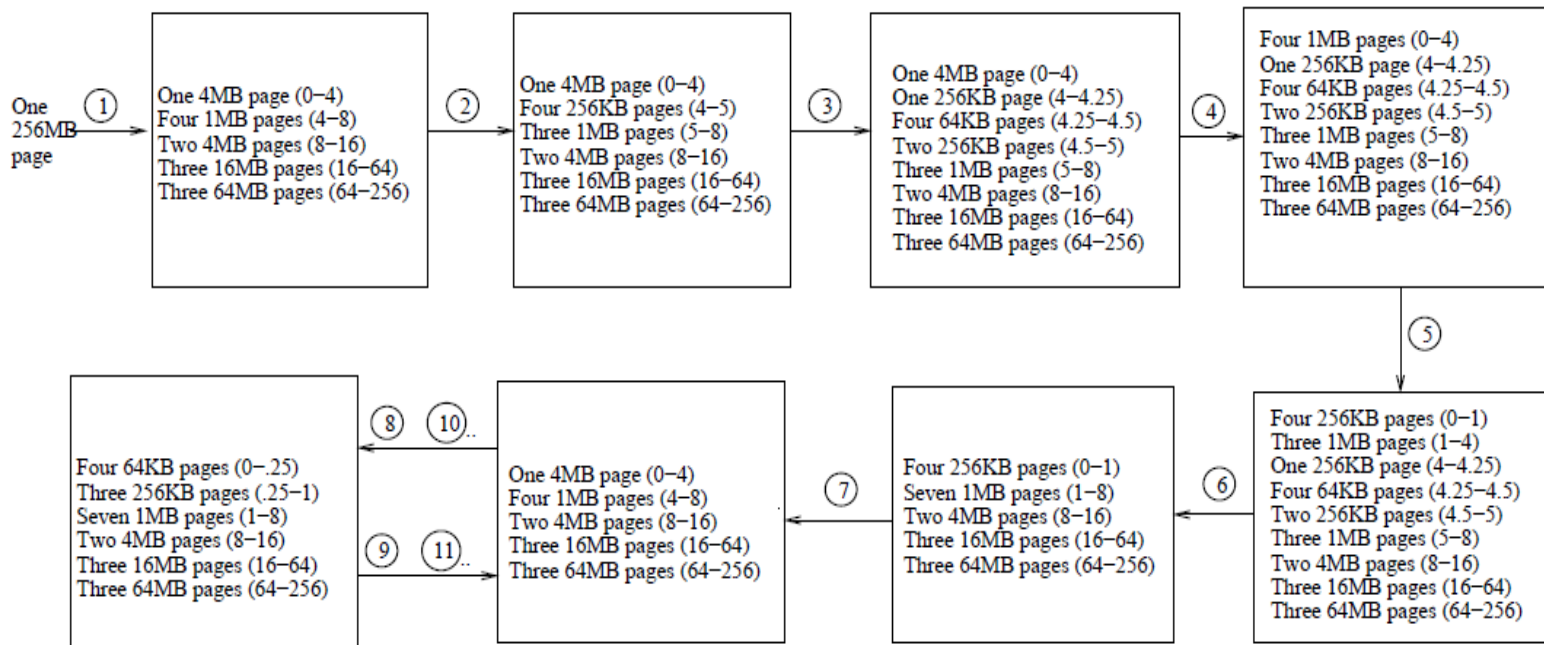
Solutions to Design Challenges



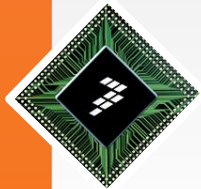
- Adaptive data mirroring algorithm to address self referential code
 - Data causing excessive DSIs mirrored to a guest memory region having R/W permission
- Addressing Translation Cache placement issue
 - Mechanism developed for stealing space from read only guest section
 - The data from the stolen section mirrored to new location



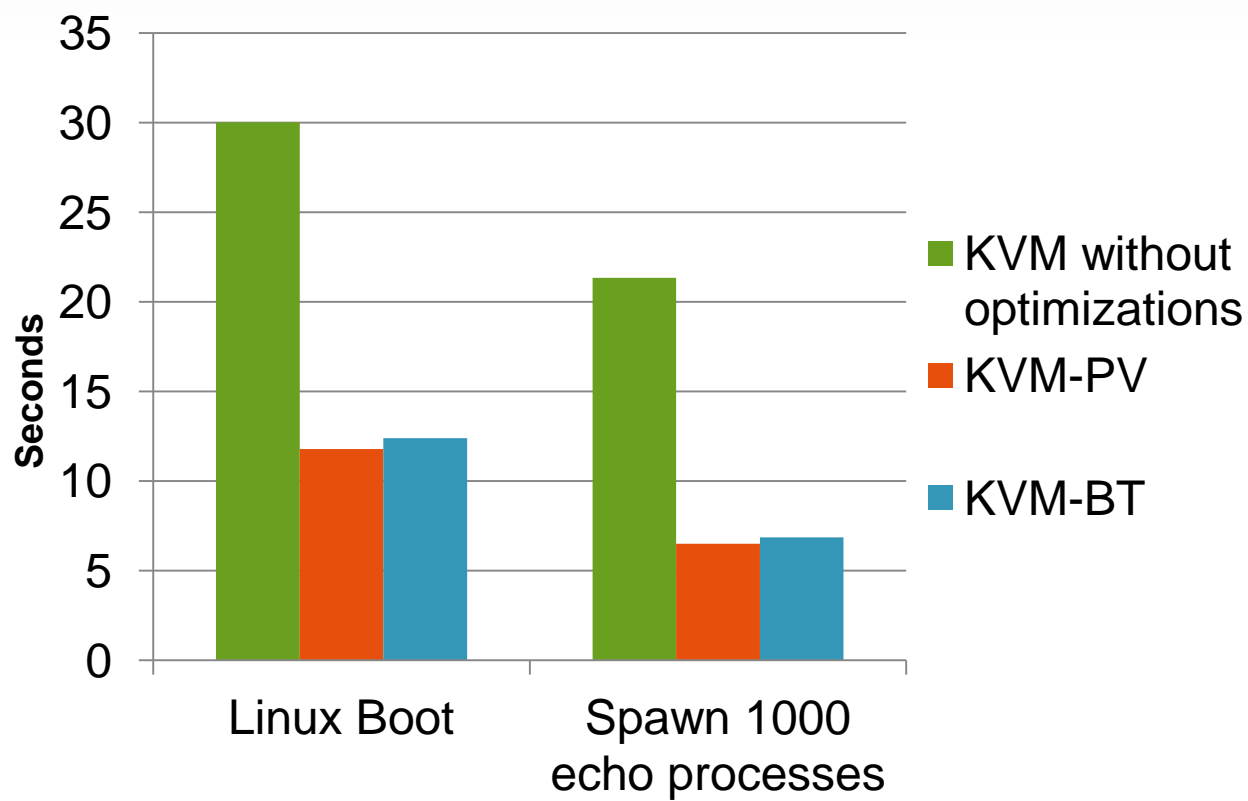
Solutions to Design Challenges

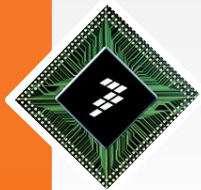


- Addressing Memory tracing issues:
 - Adaptive page resizing algorithm for addressing issue arising out of the use of Huge TLB and self modifying code
 - Dynamically splitting/merging traced pages



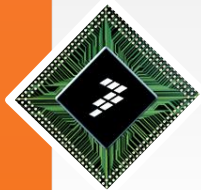
Performance





Summary

- With “Adaptive Binary Translation” it’s possible to run an unmodified guest efficiently with embedded Power Architecture KVM
- Possible to mitigate memory tracing overhead using “Adaptive Page Resizing” and “Adaptive Data Mirroring”
- Performance is comparable to the existing PV solution



Questions?

