COLO: COarse-grain LOck-stepping Virtual Machine for Non-stop Service

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Agenda

VM Replication & COLO

- COLO_KVM
- Performance Prediction
- Summary

Non-Stop Service with VM Replication

- Typical Non-stop Service Requires
 - Expensive hardware for redundancy
 - Extensive software customization
- VM Replication: Cheap Application-agnostic Solution



Existing VM Replication Approaches

- Replication Per Instruction: Lock-stepping
 - Execute in parallel for deterministic instructions
 - Lock and step for un-deterministic instructions
- Replication Per Epoch: Continuous Checkpoint
 - Secondary VM is synchronized with Primary VM per epoch
 - Output is buffered within an epoch

Problems

Lock-stepping

- Excessive replication overhead
 - memory access in an MP-guest is undeterministic
- Continuous Checkpoint
 - Extra network latency
 - Excessive VM checkpoint overhead

Why COarse-grain LOck-stepping (COLO)

VM replication is an overly strong condition

- Why do we care about the VM state ?
 - The client cares about response only
- Can the control failover without "precise VM state replication"?
- Coarse-grain lock-stepping VMs
 - Secondary VM is a replica, if it has generated same response as the primary so far
 - If true, failover with no service stop

Non-stop service focus on server response, not internal machine state!

How COLO Works

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Response Model for C/S System

 $R_n = g_n(r_0, r_1, r_2, \dots, r_n, u_0, \dots, u_m)$

- *r_i* & *u_i* are the request and the execution result of an non-deterministic instruction
- Each response packet from the equation is a semantics response

□ Successfully failover at *kth* packet if $C = \{R_1^p, ..., R_k^p, R_{k+1}^s, ...\}$ $\forall i \leq k, R_i^s = R_i^p$

(*C* is the packet series the client received)

Why is CoLo Better

- Comparing with Continuous VM checkpoint
 - No buffering-introduced latency
 - Less checkpoint frequency
 - On demand vs. periodic
- Comparing with lock-stepping
 - Eliminate excessive overhead of undeterministic instruction execution due to MP-guest memory access

Status

- Academia paper published at ACM
 Symposium on Cloud Computing (SOCC'13)
 - "COLO: COarse-grained LOck-stepping Virtual Machines for Non-stop Service"
 - <u>http://www.socc2013.org/home/program</u>
 - Refer to the paper for technical details
- Industry announcement
 - Huawei FusionSphere uses COLO
 - <u>http://enterprise.huawei.com/ilink/enenterprise/about/news/news-list/HW_308817?KeyTemps</u>

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Architecture of COLO

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Pnode: primary node; PVM: primary VM; Snode: secondary node; SVM: secondary VM

Network Process



No need to change exiting vNIC modules – use *tc* for packet redirect / mirror

RX

Pnode

- Receive a packet from client
- Mirror the packet and send to Snode
- Send the packet to Tap NIC

Snode

- Receive the packet from Pnode
- Redirect the packet to Tap NIC

ТΧ

Snode

Redirect the SVM packet to Pnode

Pnode

- Redirect the PVM packet to ifb0
- Redirect the SVM packet to ifb1
- CM compares PVM/SVM packet in ifb
 - Same: send the packet to client
 - Different: trigger checkpoint

Storage Process



Need modify Qemu vDisk IO path – intercepted by Colo Disk Manager (DM)

Write

Pnode

- DM sends the Write request (offset, len, data) to PVM cache in Snode
- DM calls block driver to write to storage

Snode

• DM saves Write request in SVM cache

Read

- Snode
 - From SVM cache, or storage otherwise

Pnode

From storage

Checkpoint

• DM calls block driver to flush PVM cache

Failover

• DM calls block driver to flush SVM cache

Storage Process (2)

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······Origin ops(skip) ······

Checkpoint Process



Need modify migration process in Qemu to support checkpoint

Checkpoint + Memory Sync Process



Need modify migration process in Qemu to support checkpoint

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VM Replication & COLO COLO_KVM Performance Prediction

Summary

Web Server Performance - Web Bench



Web Server Performance - Web Bench (MP)



Predication base on data in SOCC'13 paper

PostgreSQL Performance - Pgbench



PostgreSQL Performance - Pgbench (MP)



Predication base on data in SOCC'13 paper

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Summary

- COLO is an ideal Application-agnostic Solution for Non-stop service
 - Web server: 67% of native performance
 - CPU, memory and netperf: near-native performance
- Next steps
 - Redesign based on feedback
 - Develop and send out for review
 - Optimize performance